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# Functional Programming in

 **Haskell**

(What the hell is a monad?)

- Program is made of several steps the computer is to carry out
- You tell the computer **what** to do (imperative)
- What you're probably used to
  - Python, Java, C, JavaScript, Lua

- Program is made of functions that map values onto other values
- You tell the computer **what you want** (declarative)
  - Lisp, Erlang, Elixir, OCaml, F#, Scala

$$f(x) = x^2$$

- You have to give up lots of things you're used to
  - Mutation (`x = 1; x = 2;`)
  - Looping (use recursion instead)
  - Side effects (IO, globals and state)
    - (not really, but it's kind of funky)

- Being used more in mainstream languages
- Often results in code that is shorter and easy to read
- Allows you to do complex operations very easy
- Clout



```
reverseList :: [a] -> [a]
reverseList [] = []
reverseList (x:xs) = reverseList xs ++ x

-- λ reverseList [4,3,2,1]
-- [1,2,3,4]
-- λ reverseList "abc"
-- "cba"
```

- Haskell has a powerful type system that's clear to read
- `mult :: Int -> Int -> Int`
  - Takes two ints, returns an int
- `multBy :: Int -> (Int -> Int)`
  - Takes an int, returns a function which takes an int and returns an int

- `reverseList :: [a] -> [a]`
  - Takes a list of anything, and returns a list of anything
- `allSame :: Eq a => [a] -> Bool`
  - Takes a list of anything that is an instance of `Eq`, returns a boolean
  - Being an instance of `Eq` means two of a type can be equal (`==`) or not equal (`/=`)



- In Haskell, functions can only take on parameter
- But, functions can return functions
- Similar to idea of closures in procedural languages
- You can still use them as though they're taking multiple arguments, without having to worry about currying behind the scenes

- Allows for partial application



```
mult :: Int -> Int -> Int  
mult x y = x * y
```

```
double :: Int -> Int  
double = mult 2
```

- To Haskell, parentheses in type definitions **do not matter** unless taking functions as parameters
  - `mult :: Int -> Int -> Int`
  - `multBy :: Int -> (Int -> Int)`
  - Both of these mean the exact same thing
- But they usually help a reader better understand what your function is supposed to mean

- Haskell offers many ways to give and accept parameters
- Regular
- Guards
- Pattern matching

- Think of splitting an if-else block into separate functions




```
isOdd :: Int -> Bool
isOdd x | mod x 2 == 0 = False
        | otherwise   = True
```

- Possibilities are attempted, going down
- If none match, error is thrown
  - A side effect (gross)




```
int2str :: Int -> String
int2str 0 = "Zero"
int2str 1 = "One"
int2str 2 = "Two"
int2str 3 = "Three"
int2str 4 = "Four"
```

- Underscore discards argument



```
mult :: Int -> Int -> Int
mult 0 _ = 0
mult _ 0 = 0
mult = (*)
```

- $(x:xs)$  pattern puts first element in  $x$ , rest in  $xs$



```
sum :: [Int] -> Int
sum [] = 0
sum (x:xs) = x + sum xs

length :: [a] -> Int
length [] = 0
length (_:xs) = 1 + length xs
```



- Lists and recursion are your replacement for loops
- All elements are of same type
- Lists are constructed one element at a time, with :
- Lists can be concatenated with ++
- Random access with [1,2,3,4] !! 0



```
λ x = [1..4]
λ x
-- [1,2,3,4]

-- Get first element
λ head x
-- 1
-- Last elements
λ tail x
-- [2,3,4]

-- Take n elements
λ take 2 x
-- [1,2]
-- Drop n elements
λ drop 2 x
-- [3,4]
```



```
-- Get last element
λ last x
-- 4
-- Drop last element
λ init x
-- [1,2,3]

-- Check if list is empty
λ null x
-- False

-- Check if element in list
λ elem 3 x
-- True
-- P.S. You can use functions with two args
-- like operators by doing this:
λ 5 `elem` x
-- False
```

- Generate a new list from an existing one very easily

```
λ [int2str i | i <- [1..5]]
-- ["One","Two","Three","Four","Five"]
λ [int2str i | i <- [1..10], i `mod` 2 == 0]
-- ["Two","Four","Six","Eight","Ten"]
λ [(int2str x, int2str y) | x <- [1..3], y <- [4..5]]
-- [("One","Four"),
--  ("One","Five"),
--  ("Two","Four"),
--  ("Two","Five"),
--  ("Three","Four"),
--  ("Three","Five")]
```



```
-- Map a function over a list
λ map int2str [1..4]
-- ["One","Two","Three","Four"]

-- Type of [1..4] is [Int]
-- Type of int2str is Int -> String
-- Type of output is [String]
-- Type of map is (a -> b) -> [a] -> [b]
```



```
-- Filter a list with a predicate
λ filter isOdd [1..4]
-- [1..3]
-- Type of filter is (a -> Bool) -> [a] -> [a]

-- Combine two lists into tuples
λ zip [1..3] [4..6]
-- [(1,4),(2,5),(3,6)]
```

- Quick and dirty functions, ideal for map and filter



```
λ (\a -> a + 2) 1  
-- 3
```

```
λ (\a b -> a + b) 3 5  
-- 8
```

```
λ filter (\a -> isOdd a && a > 4) [1..10]  
-- [5,7,9]
```

- `(. )` operator works similarly to function composition from math

```
halve :: Int -> String
halve = int2str . (\x -> div x 2)

-- Type of (.) is (b -> c) -> (a -> b) -> (a -> c)
-- so it can't work on functions with more than one arg.
-- You can take advantage of currying and do this:

myDiv :: Int -> Int -> String
myDiv = (int2str . ) . div

-- But it's probably more readable to do it the more
-- obvious way:

myDiv x y = int2str (div x y)
```

- Many recursive functions follow a common pattern:
  - Base case
  - Operation that takes next element and accumulation

- When applied to [1..4]:
  - $(1 + (2 + (3 + (4 + 0))))$
  - $(1 + (2 + (3 + (4))))$
  - $(1 + (2 + (7)))$
  - $(1 + (9))$
  - $(10) \rightarrow 10$



```
sum :: [Int] -> Int
-- foldr takes a function, and a base case
sum = foldr (\x acc -> x + acc) 0
```



- Order of ops reversed (folds left)
- 1 variant uses first element as base case



```
subList :: [Int] -> Int
subList = foldl1 (\acc x -> acc - x)

λ subList [10,7,2]
-- 1
```

- Structuring your own data is easy, and there are several tools for you
  - Types
  - Data
  - Typeclasses and Newtypes

- Just an alias for an existing type, good for readability



```
type Pair a = (a, a)
type Pos = Pair Int
```

```
-- Haskell will consider all of these to be
-- the same type
-- (5, 5)
-- (4, 3) :: Pair Int
-- (-1, 2) :: Pos
```

- Similar to a struct in C, with some new tricks

```
data Person = PersonPhone String String Int | PersonEmail String String String
    deriving Show
-- deriving Show automatically makes it an instance of the Show typeclass, so it can
-- be printed to the console

λ PersonPhone "John" "Doe" 1234567890
-- PersonPhone "John" "Doe" 1234567890
-- It works, but you have to remember what goes where,
-- and to use it in functions you have to do this:
-- f (PersonPhone fstName lstName phone) = ...

data BetterPerson = BPersonPhone { bppFirstName :: String
                                   , bppLastName  :: String
                                   , bppPhone      :: Int }
    | BPersonEmail { bpeFirstName :: String
                    , bpeLastName  :: String
                    , bpeEmail     :: String }

    deriving Show

-- Now you can do this
λ jdoe = BPersonEmail { bpeFirstName = "John", bpeLastName = "Doe", bpeEmail = "jdoe@gmail.com"}
λ bpeFirstName jdoe
-- "John"
```



```
data NaturalNumber = Zero | Succ NaturalNumber
    deriving Show
```

```
nat2int :: NaturalNumber -> Int
nat2int Zero = 0
nat2int (Succ n) = 1 + nat2int (n-1)
```

```
int2nat :: Int -> NaturalNumber
int2nat 0 = Zero
int2nat n = Succ (int2nat (n-1))
```

```
λ int2nat 3
-- Succ (Succ (Succ Zero))
λ nat2int (Succ (Succ Zero))
-- 2
```

- Defines a set of functions for a given type

```
class Eq a where
  (==) :: a -> a -> Bool
  (/=) :: a -> a -> Bool
  x == y = not (x /= y)
  x /= y = not (x == y)

instance Eq NaturalNumber where
  Zero == Zero = True
  Zero == _ = False
  _ == Zero = False
  (Succ n1) == (Succ n2) = nat2int n1 == nat2int n2
```

- Similar to data, but only allows for one constructor and field
- Used to wipe typeclass definitions for an existing type
- Also comes with some performance benefits



```
newtype MyList a = MyList { getList :: [a] }  
  
-- Can't use !!, :, or ++ with MyLists  
-- We're free to define them ourselves now
```

- Functors 😊
- Applicatives 😐



- Monads




- Easy way to model a computation that may fail



```
data Maybe a = Just a | Nothing
    deriving (Eq, Ord)
```

```
safeDiv :: Int -> Int -> Maybe Int
safeDiv _ 0 = Nothing
safeDiv x y = Just (x `div` y)
```

- With a functor, we can map a function onto its inner value(s)



```
instance Functor Maybe where
    fmap f (Just x) = Just (f x)
    fmap f Nothing = Nothing

λ fmap (+2) (safediv 9 3)
-- Just 5
λ fmap (+2) (safediv 9 0)
-- Nothing
```

- map is fmap for lists




```
instance Functor [] where
    fmap _ [] = []
    fmap f (x:xs) = f x : fmap f xs
```

- Functors are cool, but they only work on functions with one argument
- Applicatives solve this problem
- They allow you to wrap a function within a type (pure), and then feed it parameters

# Maybe as an applicative

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```
instance Applicative Maybe where
    pure = Just
    Nothing <*> _ = Nothing
    (Just f) <*> something = fmap f something

λ pure (+) <*> Just 2 <*> Just 3
-- Just 5

-- A little bit of sugar
λ (*) <$> Just 3 <*> Just 4
-- Just 12
λ (\a b c -> a * b + c) <$> Just 10 <*> Just 5 <*> Just 25
-- Just 75

λ (+) <$> Nothing <*> Just 4
-- Nothing
λ (*) <$> Just 3 <*> Nothing
-- Nothing
```

- We have a process that could fail at any step
- Let's use Maybe
- Each step needs to be able to pass on failure (Nothing) from the step before





- Can we use functors?
  - No, because that only applies the function to the inner value. If the function fails, we'd still get an exception, or at best, `Just Nothing`
- What about applicatives?
  - No, because we have to start with a function and feed values into it. We can't chain the output of one function to the output of another

- We need to be able to:
  - Chain together as many functions as we want
  - Be able to propagate failure to the end of the chain, from anywhere in the chain
    - Each function needs to be able to return `Nothing`
    - But we also want our functions to be as generic as possible, so their parameters shouldn't be wrapped in `Maybe`
    - Their types should be `a -> Just a`
    - But now, how do we pass output (`Just a`) to the next function (which expects an `a`)?

- Let's make an intermediary function to handle it

```
bind :: Maybe a -> (a -> Maybe a) -> Maybe a
bind Nothing _ = Nothing
bind (Just x) f = f x

safeDiv :: Int -> Int -> Maybe Int
safeDiv _ 0 = Nothing
safeDiv x y = Just (x `div` y)

divBy :: Int -> Int -> Maybe Int
divBy x y = safeDiv y x

λ safeDiv 24 2 `bind` divBy 2 `bind` divBy 3
-- Just 2
λ safeDiv 24 0 `bind` divBy 2 `bind` divBy 3
-- Nothing
λ safeDiv 24 2 `bind` divBy 0 `bind` divBy 3
-- Nothing
λ safeDiv 24 2 `bind` divBy 2 `bind` divBy 0
-- Nothing
```

- This is really handy, someone should name this!
- monad
- Oh
- But hey, that wasn't so bad, was it?

- In practice, `bind` is replaced with `>>=`
- And there are many different monads, `Maybe` is just one of them
- But they all follow a similar idea, of being able to pass failure, or really any context, down a pipeline

- Haskell functions are pure, they have no side effects
- ... but every way we could interact with a program is a side effect
- So is Haskell actually pure?



- *Every* Haskell function *is* pure
- So to get around it, we can make a type called IO, which is a side effect
- Functions can now return side effects, **but the function itself is not what causes a side effect to happen**

- To execute a side effect, set `main` equal to it



```
main = putStrLn "Hello, World!"
```

```
λ :type main
```

```
-- IO ()
```

```
-- IO types can also wrap a value
```

```
λ :type getLine
```

```
-- IO String
```

```
main = getLine >=> (\s -> putStrLn ("Hello, " ++ s))
```



- The flow of your program can be represented purely, so you're not really breaking any "purely functional" rules
- Haskell just executes the IO action you've given it, and uses the rules you've defined in your functions

- You'll probably want to do more than one or two things in your main action
- This sucks

```
main = putStrLn "First name?" >>
      getLine >>= (\firstName ->
        putStrLn "Last name?" >>
        getLine >>= (\lastName ->
          putStrLn "Email?" >>
          getLine >>= (\email ->
            putStrLn (lastName ++ ", " ++ firstName ++ " (" ++ email ++ ")"))))
```

- Haskell offers some more sugar
- This translates straight to what was seen earlier, just easier to read
- In fact, do syntax works on any monad, not just IO

```
contactStr :: String -> String -> String -> String
contactStr f l e = l ++ ", " ++ f ++ " (" ++ e ++ ")"

main = do
    putStrLn "First name?"
    -- bind result of IO to a variable
    firstName <- getLine
    putStrLn "Last name?"
    lastName <- getLine
    putStrLn "Email?"
    email <- getLine
    let result = contactStr firstName lastName email
    -- could also do:
    -- result <- return (contactStr firstName lastName email)
    -- which wraps the String from contactStr in IO
    -- but there's no point to it since we're binding it right after
    putStrLn result
```

- Since mutability is not supported, a function will always return the same output for the same input
  - Don't worry about IO for now
- Because of this, an expression will always evaluate the same
  - Also makes parallelization/concurrency very easy
- Haskell is lazy: it doesn't evaluate an expression until it absolutely needs to



```
hardMathProblem :: Int
hardMathProblem = (code to actually calculate it)
```

```
λ let x = hardMathProblem
-- Interpreter sets x equal to hardMathProblem. No calculation
-- is actually performed because Haskell doesn't have to show
-- you the result yet.
λ x
-- Now, Haskell needs to figure out what x really is. This is
-- where computation starts, and where the interpreter would pause.
λ 3.14159.....
```

- This has performance benefits (sometimes), since unnecessary computation isn't performed
  - Also makes debugging harder :(
- But it also allows for one very cool trick that is impossible (or at least much less natural) in a non-lazy language



```
λ x = [1..]
-- x is now the list of all integers, starting at 1
λ take 5 x
-- [1,2,3,4,5]

-- Of course, this doesn't work, but Haskell will still try
-- sum :: [Int] -> Int
λ sum x

-- You can also make your own infinite data structures.
-- This hurts my head, but Haskell doesn't blink.
primes = filterPrime[2..]
  where filterPrime (p:xs) =
    p : filterPrime [x | x <- xs, x `mod` p /= 0]

λ take 5 primes
-- [2,3,5,7,11]
```



```
-- Let's say you have a config file on your drive, and you have a function
-- to open and parse it into a type Config:
let conf = parseConfigFile "config.json"
-- conf's type is IO Config
-- Now let's say you try to use it, but since typing the above line, you've
-- deleted the config file.
λ conf
-- Error! File not found
-- Haskell didn't try to open the file until it really needed to

-- However, by instead binding the result of parseConfigFile:
conf <- parseConfigFile "config.json"
-- IO Config is unwrapped into just Config, and Haskell is forced to get its
-- true value even if you don't necessarily need it yet.
```




- Glasgow Haskell Compiler (GHC)
  - The GCC of Haskell
  - Can also interpret, run a REPL
- Hackage: Package index
- Stack: Build system, package manager
- Cabal: Older build system
  - Still updated, but falling out of favor for Stack
- Haddock: Documentation generator
  - Like JavaDoc, Doxygen

- Learn You a Haskell For Great Good!
  - Miran Lipovaca
  - Available for free online
- Programming in Haskell
  - Graham Hutton (of Computerphile fame)
- Real World Haskell
  - O'Reilly
- Tsoding on YouTube

Some examples were taken from LYAHFGG, Programming in Haskell, and the Haskell website.

`carbon.now.sh` was used to generate code snippets.



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